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**A FORTRAN PROGRAM FOR THE CALCULATION OF THE
STATE TRANSITION MATRIX AS A LINEAR COMBINATION
OF REAL TIME FUNCTIONS (EAT)**

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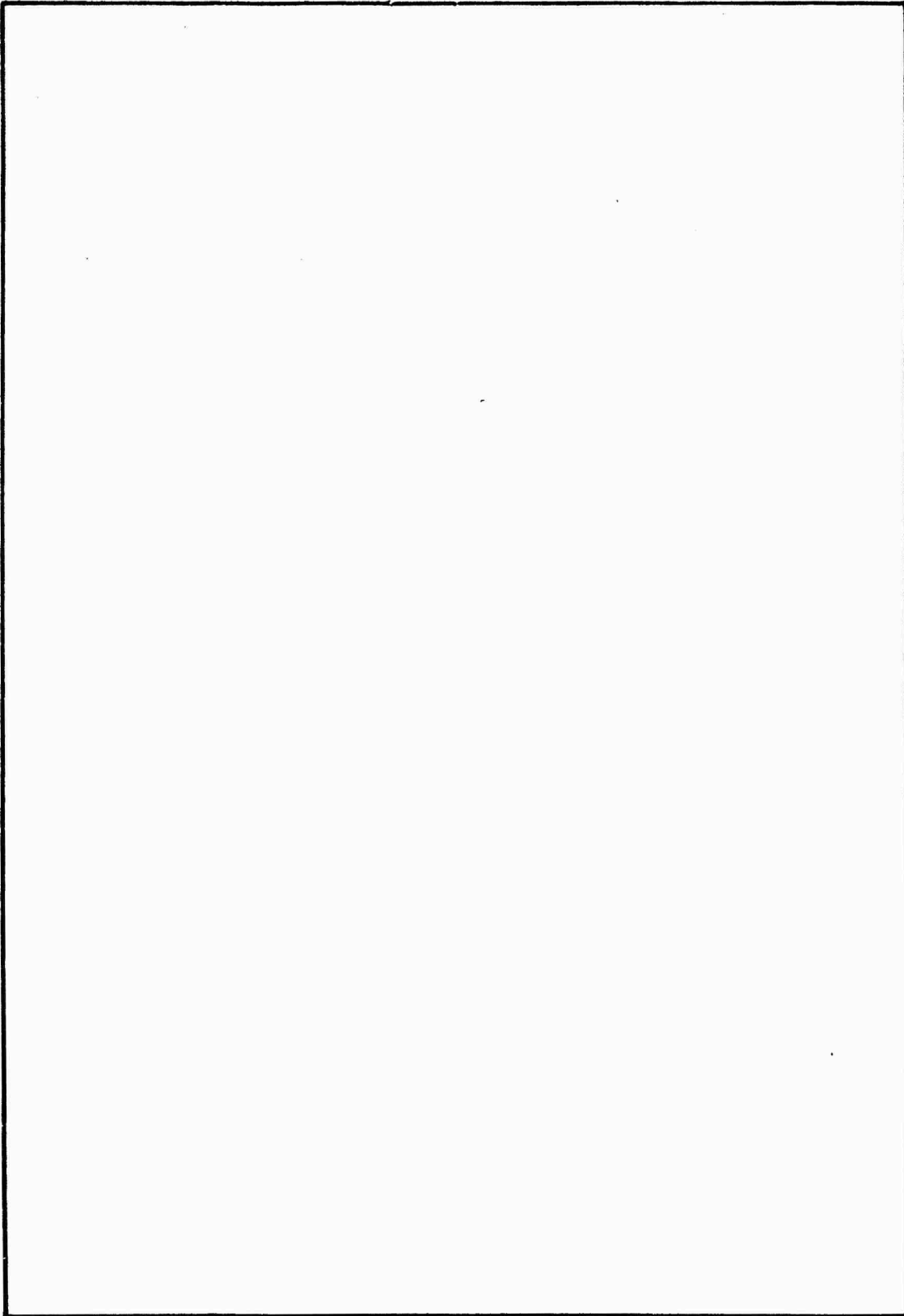
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20. ABSTRACT (Continue on reverse side if necessary and identify by block number) Many programs are available which calculate numerically the state transition matrix at a specific time. However, programs which give the solution as a linear combination of real time functions are not generally available. The program given in this report calculates the state transition matrix as a linear combination of real time functions starting from any real system matrix.		

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1. Introduction

The authors have developed a FORTRAN program that can be used to calculate the solution to the homogeneous system of first order differential equations with constant coefficients as follows:*

$$\begin{aligned}\dot{\bar{Y}}(t) &= \underline{A}\bar{Y}(t) \\ \bar{Y}(t_0) &= \bar{Y}_0\end{aligned}\tag{1}$$

Consider Equation (1) and let $\bar{Y}(t)$ be an n -vector of differentiable functions, \bar{Y}_0 be an n -vector of constants (initial conditions), and \underline{A} be an $n \times n$ matrix with constant elements. Then the solution to Equation (1) can be written as

$$\bar{Y}(t) = \underline{\Phi}(t - t_0)\bar{Y}_0\tag{2}$$

where $\underline{\Phi}$ is an $n \times n$ matrix known as the state transition matrix (fundamental matrix) whose entries are functions of t . For example,

$$\begin{aligned}\dot{y}_1(t) &= y_2(t) \\ \dot{y}_2(t) &= -2y_1(t) - 2y_2(t)\end{aligned}\tag{3}$$

is a 2×2 system of linear first order equations, and $\underline{\Phi}(t)$ is the matrix

$$\begin{pmatrix} e^{-t}(\cos t - \sin t) & e^{-t} \sin t \\ -2 e^{-t} \sin t & e^{-t}(\cos t - \sin t) \end{pmatrix}.\tag{4}$$

*The symbol \underline{A} will be used to indicate A is a matrix and the symbol \bar{A} will be used to indicate A is a column vector.

To determine a solution to the system, given a set of initial conditions,

$$\bar{Y}(t_0) = \begin{pmatrix} y_1(t_0) \\ y_2(t_0) \end{pmatrix} = \bar{Y}_0 = \begin{pmatrix} y_1 \\ y_2 \end{pmatrix}, \quad (5)$$

one need only premultiply it by $\Phi(t - t_0)$. This result has been in the literature for quite sometime¹ and is a special case of the more general result that

$$\begin{aligned} \bar{Y}(t) &= \underline{A}\bar{Y}(t) + \underline{B}\bar{V}(t) \\ \bar{Y}(t_0) &= \bar{Y}_0 \end{aligned} \quad (6)$$

is solved by

$$\Phi(t - t_0)\bar{Y}_0 + \int_{t_0}^t \Phi(t - t_0)\underline{B}\bar{V}(t)dt. \quad (7)$$

Clearly, the central problem in determining a particular solution in either the homogeneous or the more general problem is the calculation of $\Phi(t)$, better known as $e^{\underline{A}t}$. Subroutines for calculating $e^{\underline{A}t}$ have been in the literature for quite sometime. All those known to the authors, however, have the drawback that they do not output $e^{\underline{A}t}$ in a form similar to that of Equation (4), but give $e^{\underline{A}t}$ for one value of $t = t_a$. This technique is widely used in the numerical integration of linear systems. Though such techniques are useful, the analytical form of the solution is lost together with time constants and system frequencies which appear in an analytical representation for $\Phi(t)$.

The computer program developed in this report can be used to obtain $\Phi(t)$ for every t . The user need only input the system matrix \underline{A} , the dimension of \underline{A} , and a set of error tolerance levels. For example, let

$$\underline{A} = \begin{pmatrix} 1 & 0 \\ -2 & -2 \end{pmatrix}. \quad (8)$$

¹Frame, J. S., "Matrix Functions and Applications, IV," IEEE Spectrum, June 1964, pp. 123-131.

Then the program output will contain (among other things) $\underline{\Phi}(t)$ expressed as a sum of matrices times linearly independent functions:

$$\underline{\Phi}(t) = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} e^{-t} \cos t + \begin{pmatrix} -1 & 1 \\ -2 & -1 \end{pmatrix} e^{-t} \sin t \quad (9)$$

The error tolerance levels are discussed in detail in Section 3.

The remainder of this report is divided into three sections. Section 2 gives a description of the problem and the techniques used to solve it. Section 3 deals exclusively with the computer program and includes a description of inputs and outputs. Finally, Section 4 gives a listing of the program. Those familiar with analytic functions of matrices can go directly to Section 3.

2. The State Transition Matrix

a. $e^{\underline{A}t}$, Definition, and Properties

$e^{\underline{A}t}$, the $\underline{\Phi}(t)$ of Section 1, is defined by a power series as

$$e^{\underline{A}t} = \underline{I} + \underline{A}t + \frac{\underline{A}^2 t^2}{2!} + \frac{\underline{A}^3 t^3}{3!} + \dots \quad (10)$$

where \underline{I} represents the unit matrix. Since $e^{\underline{A}t}$ is defined analogously to $e^{\alpha t}$, α is a scalar, it is similar in many respects. For example,

$$\begin{aligned} e^{\underline{A}(t+s)} &= e^{\underline{A}t} e^{\underline{A}s} \\ \frac{d}{dt} e^{\underline{A}t} &= \underline{A} e^{\underline{A}t} \\ e^{\underline{A} \cdot 0} &= \underline{I} \quad ; \end{aligned} \quad (11)$$

but, generally speaking

$$\begin{aligned} e^{(\underline{A}t + \underline{B}t)} &\neq e^{\underline{A}t} e^{\underline{B}t} \\ e^{\underline{A}t} e^{\underline{B}t} &\neq e^{\underline{B}t} e^{\underline{A}t} \end{aligned} \quad (12)$$

since matrix multiplication is not a commutative operation. A computationally useful property of $e^{\underline{A}t}$ is

$$e^{\underline{R} \underline{A} \underline{R}^{-1} t} = \underline{R} e^{\underline{A} t} \underline{R}^{-1} \quad (13)$$

This makes it possible to easily calculate $e^{\underline{A} t}$ in some cases by transforming it into a similar matrix (e.g., normal matrices to diagonal matrices²). The power series representation, while useful for numerical work, is none too helpful for writing $e^{\underline{A} t}$ in closed form. The following general theorem is used to decompose $e^{\underline{A} t}$ into a finite sum of $n \times n$ matrices times analytic functions in one variable.

Theorem³

If f is an analytic function on a simply connected open set D of the complex plane that contains all the eigenvalues λ_j of an $n \times n$ matrix \underline{B} and the origin, then $f(\underline{B})$ may be written as

$$f(\underline{B}) = \sum_{j=1}^s \sum_{k=1}^{n_j} \frac{f^{(k-1)}(\lambda_j)}{(k-1)!} \underline{Z}_{j,k} \quad (14)$$

where n_j denotes the multiplicity of the j^{th} eigenvalue, s is the number of eigenvalues, and the $\underline{Z}_{j,k}$ are $n \times n$ matrices which are independent of f and D (they depend only on the matrix \underline{B}).

Since e^x is analytic in the whole complex plane, setting $\underline{B} = \underline{A} t$, $f(x) = e^x$ yields

$$e^{\underline{A} t} = \sum_{j=1}^s \sum_{k=1}^{n_j} \frac{t^{k-1} e^{\lambda_j t}}{(k-1)!} \underline{Z}_{j,k} \quad (15)$$

Now it is clear that $e^{\underline{A} t}$ has a representation as a finite sum of $n \times n$ matrices of scalars times analytic functions. The problem is to determine the λ_j and the $\underline{Z}_{j,k}$.

²Herstein, I. N., Topics in Algebra, Blaisdell, Waltham, Massachusetts, 1964.

³Frame, loc. cit.

A review of the terminology involved in the theorem is presented. A function analytic about $\{0\}$ has a power series representation about $\{0\}$:

$$f(z) = \sum_{n=0}^{\infty} a_n z^n \text{ for } |z| < r ; \quad (16)$$

therefore, $f(\underline{B})$ can be formally defined by

$$f(\underline{B}) = \sum_{n=0}^{\infty} a_n \underline{B}^n . \quad (17)$$

If \underline{B} is an $n \times n$ matrix, an eigenvalue of \underline{B} is a scalar λ_j for which a vector \bar{x} exists such that

$$\underline{B}\bar{x} = \lambda_j \bar{x} \quad (18)$$

or

$$(\underline{B} - \lambda_j \underline{I})\bar{x} = \underline{0} . \quad (19)$$

If Equation (19) holds, the matrix $\underline{B} - \lambda_j \underline{I}$ is singular, and, therefore, λ_j is a solution of the equation,

$$\text{Det}(\underline{B} - \lambda_j \underline{I}) = 0 , \quad (20)$$

where $\text{Det}(\underline{B} - \lambda_j \underline{I})$ is an n^{th} degree polynomial called the characteristic polynomial of \underline{B} . Its roots, which are just the eigenvalues of \underline{B} , are also called the characteristic roots of \underline{B} . The multiplicity of λ_j is its multiplicity as a root of $\text{Det}(\underline{B} - \lambda_j \underline{I})$.

The theorem asserts that $f(\underline{B})$ exists (the power series of Equation (17) converges) if the eigenvalues are in D and that $f(\underline{B})$ has a representation in the form of Equation (14).

b. The Characteristic Polynomial and Its Roots

Certainly, the first problem in writing $e^{\underline{A}t}$ in the form of Equation (15) is to calculate the eigenvalues and determine their multiplicities. To do this the characteristic polynomial has to be calculated. If

$$\text{Det}(\underline{A} - \lambda \underline{I}) = \lambda^n + d_1 \lambda^{n-1} + \dots + d_0 \quad (21)$$

is the characteristic polynomial, the coefficients d_k are calculated using the following theorem.

Theorem⁴

$$d_k = -\text{tr}(\underline{A}\underline{B}_{k-1})$$

$$\underline{B}_k = \underline{A}\underline{B}_{k-1} + d_k \underline{I}$$

$$d_0 = 1, \underline{B}_0 = \underline{I}, \underline{B}_n = \underline{0}$$

$$0 \leq k \leq n \quad . \quad (22)$$

Recall that $\text{tr}(a_{ij}) = \sum a_{ii}$ is the sum of the diagonal entries of (a_{ij}) . Since $\underline{B}_n = \underline{0}$ a good way to check for errors in the calculation of the coefficients is to check the difference between the calculated \underline{B}_n and the theoretically determined values. An interesting consequence of Equation (21) is that

$$\frac{\underline{A}\underline{B}_{n-1}}{-d_n} = \underline{I} \quad (23)$$

or

$$\frac{\underline{B}_{n-1}}{-d_n} = \underline{A}^{-1} \quad , \quad (24)$$

if \underline{A}^{-1} exists ($d_n \neq 0$).

The roots of $\text{Det}(\underline{A} - \lambda \underline{I})$ can, at this stage, be calculated using any one of a number of different techniques. In this program the classical Newton-Raphson method is used. The roots then have to be sorted and multiplicities counted. Numerical errors can be generated almost anywhere in our program. At this stage, these errors necessitate a decision. For example, using Newton-Raphson the equation

$$x^4 + 2x^2 + 1 = 0 \quad (25)$$

will have calculated roots $\epsilon_1 \pm i$, $\epsilon_2 \pm i$ where ϵ_1, ϵ_2 are small ($< 10^{-13}$ using a double precision version of an IBM root routine), but distinct real numbers. The solution to a 4×4 system with Equation (25) as its characteristic equation will be calculated to be of the form

⁴Frame, op. cit.

$$\underline{z}_{1,1}e^{(\epsilon_1+i)t} + \underline{z}_{2,1}e^{(\epsilon_1-i)t} + \underline{z}_{3,1}e^{(\epsilon_2+i)t} + \underline{z}_{1,4}e^{(\epsilon_2-i)t} . \quad (26)$$

But since the roots of Equation (25) are really $\pm i$, each of multiplicity two, the solution really is

$$\underline{z}_{1,1}e^{it} + \underline{z}_{1,2}te^{it} + \underline{z}_{2,1}e^{-it} + \underline{z}_{2,2}te^{-it} . \quad (27)$$

To avoid this kind of problem some decision has to be made. Namely, a tolerance ϵ is specified such that a will be set equal to b if

$$|a - b| < \epsilon \quad (28)$$

where a and b are roots of the characteristic polynomial.

For example, in the case previously considered, ϵ is set to be large enough so that

$$|\epsilon_1 - \epsilon_2| < \epsilon . \quad (29)$$

The program replaces ϵ_2 by ϵ_1 and, rather than saying that the characteristic polynomial has distinct roots $\epsilon_1 \pm i$ and $\epsilon_2 \pm i$, claims that it has a root of $\epsilon_1 \pm i$ of multiplicity two. The program then checks to determine if the root, or its real or imaginary part, is small enough to be called zero, i.e.,

$$|a| < \epsilon . \quad (30)$$

If Equation (30) is satisfied, a is set equal to zero. In the example, the final output is two roots, $\pm i$ (each of multiplicity two), if ϵ is sufficiently large.

Although an error at this stage will make the solution to the system look radically different, it will probably not change any of the usual system constants in a discontinuous manner. In fact, the solution to Equation (1) depends continuously on the eigenvalues. To see this, notice from Equation (10) that the solution depends continuously on \underline{A} . Using Equation (13) it may be assumed that \underline{A} is in lower triangular form⁵

⁵Herstein, loc. cit.

$$\underline{A} = \begin{pmatrix} \lambda_1 & & 0 \\ & \ddots & \\ * & & \lambda_s \end{pmatrix} \quad (31)$$

The diagonal entries of \underline{A} are just the eigenvalues. If the roots are changed by a small amount ϵ_i , then the system is changed to one with matrix \underline{A}' where

$$\underline{A}' = \underline{A} + \underline{\epsilon} \quad (32)$$

$$\underline{\epsilon} = \begin{pmatrix} \epsilon_1 & & 0 \\ & \ddots & \\ 0 & & \epsilon_s \end{pmatrix}$$

Hence,

$$e^{\underline{A}'t} = e^{(\underline{A} + \underline{\epsilon})t} \quad (33)$$

which is continuous in the ϵ_i .

c. The Constituent Matrices

The matrices $\underline{Z}_{j,k}$ are called the constituent matrices and, as was previously noted, are dependent only on \underline{A} , not on the analytic function at which \underline{A} is evaluated. Rather than launch into a general description of the technique used to calculate the constituent matrices, they will first be calculated for a particular example and the illustrated technique generalized. Suppose

$$\underline{A} = \begin{pmatrix} 0 & 1 & 3 \\ 6 & 0 & 2 \\ -5 & 2 & 4 \end{pmatrix} \quad (34)$$

with characteristic polynomial

$$(x - 1)^2(x - 2) \quad (35)$$

which has roots $\lambda_1 = 1$ (of multiplicity two) and $\lambda_2 = 2$ (with multiplicity one). $e^{\underline{A}t}$ must be of the form

$$e^{\underline{A}t} = e^t \underline{Z}_{1,1} + te^t \underline{Z}_{1,2} + e^{2t} \underline{Z}_{2,1} \quad (36)$$

Applying Theorem 1 to the analytic functions $f(x) = x^0$, $g(x) = x$, and $h(x) = x^2$ and substituting \underline{A} for x , the resulting equations are

$$\begin{aligned}\underline{I} &= 1 \cdot \underline{Z}_{1,1} + 0 \underline{Z}_{1,2} + 1 \underline{Z}_{2,1} \\ \underline{A} &= 1 \cdot \underline{Z}_{1,1} + 1 \cdot \underline{Z}_{1,2} + 2 \cdot \underline{Z}_{2,1} \\ \underline{A}^2 &= 1 \cdot \underline{Z}_{1,1} + 2 \cdot \underline{Z}_{1,2} + 4 \cdot \underline{Z}_{2,1}\end{aligned}\quad (37)$$

By using matrix notation and considering \underline{I} , \underline{A} , \underline{A}^2 , $\underline{Z}_{j,k}$ as formal symbols Equation (37) can be written as

$$\begin{pmatrix} \underline{I} \\ \underline{A} \\ \underline{A}^2 \end{pmatrix} = \begin{pmatrix} 1 & 0 & 1 \\ 1 & 1 & 2 \\ 1 & 2 & 4 \end{pmatrix} \begin{pmatrix} \underline{Z}_{1,1} \\ \underline{Z}_{1,2} \\ \underline{Z}_{2,1} \end{pmatrix} \quad (38)$$

The matrix

$$\begin{pmatrix} 1 & 0 & 1 \\ 1 & 1 & 2 \\ 1 & 2 & 4 \end{pmatrix} \quad (39)$$

is invertible with inverse

$$\begin{pmatrix} 0 & 2 & -1 \\ -2 & 3 & -1 \\ 1 & -2 & 1 \end{pmatrix} \quad (40)$$

so

$$\begin{pmatrix} \underline{Z}_{1,1} \\ \underline{Z}_{1,2} \\ \underline{Z}_{2,1} \end{pmatrix} = \begin{pmatrix} 0 & 2 & -1 \\ -2 & 3 & -1 \\ 1 & -2 & 1 \end{pmatrix} \begin{pmatrix} \underline{I} \\ \underline{A} \\ \underline{A}^2 \end{pmatrix} \quad (41)$$

or, in equation form,

$$\underline{Z}_{1,1} = 2\underline{A} - \underline{A}^2$$

$$\begin{aligned} \underline{Z}_{1,2} &= -2\underline{I} + 3\underline{A} - \underline{A}^2 \\ \underline{Z}_{2,1} &= \underline{I} - 2\underline{A} + \underline{A}^2 \end{aligned} \quad (42)$$

The \underline{Z} 's can be easily calculated from Equation (42). In general, if \underline{A} is an $n \times n$ matrix, the column vectors

$$\overline{\underline{A}}_a = \begin{pmatrix} \underline{I} \\ \underline{A} \\ \vdots \\ \underline{A}^{n-1} \end{pmatrix} \quad (43)$$

$$\overline{\underline{Z}}_a = \begin{pmatrix} \underline{Z}_{1,1} \\ \underline{Z}_{1,2} \\ \vdots \\ \underline{Z}_{s,n_s} \end{pmatrix} \quad (44)$$

are formed and yield the matrix equation

$$\overline{\underline{A}}_a = \underline{V}^t \overline{\underline{Z}}_a \quad (45)$$

where \underline{V}^t is the transpose of an $n \times n$ matrix \underline{V} , called the Vandermonde of Equation (1). The simplest way to calculate the entries of \underline{V}^t is as follows. Partition \underline{V}^t into s , $n \times n_s$ matrices. The k^{th} matrix will be filled with entries calculated from the k^{th} eigenvalue. The entries of each of these matrices is calculated using the algorithm

$$\begin{aligned} v_{ij} &= 0 & i > j \\ v_{ij} &= 1 & i = j \\ v_{ij} &= v_{i-1,j-1} + \lambda_k v_{i-1,j} & i < j \end{aligned} \quad (46)$$

As an example, suppose \underline{A} has 3 eigenvalues λ_1 of multiplicity 3, λ_2 of multiplicity 2, and λ_3 with multiplicity 1, then \underline{V}^t is

$$\begin{pmatrix} 1 & 0 & 0 & 1 & 0 & 1 \\ \lambda_1 & 1 & 0 & \lambda_2 & 1 & \lambda_3 \\ \lambda_1^2 & 2\lambda_1 & 1 & \lambda_2^2 & 2\lambda_2 & \lambda_3^2 \\ \lambda_1^3 & 3\lambda_1^2 & 3\lambda_1 & \lambda_2^3 & 3\lambda_2^2 & \lambda_3^3 \\ \lambda_1^4 & 4\lambda_1^3 & 6\lambda_1^2 & \lambda_2^4 & 4\lambda_2^3 & \lambda_3^4 \\ \lambda_1^5 & 5\lambda_1^4 & 10\lambda_1^3 & \lambda_2^5 & 5\lambda_2^4 & \lambda_3^5 \end{pmatrix} \quad (47)$$

where \underline{V}^t is always invertible⁶ so Equation (45) can always be solved.

d. Complex Eigenvalues

In all previous examples the eigenvalues have been real. In general, the eigenvalues may be complex and the result is that λ_j and $\underline{Z}_{i,j}$ may in general be complex. The program developed in this report handles all complex computations with real computation and no FORTRAN complex declarations are made. As a result, double precision may be used to provide accurate solutions.

If the eigenvalues are complex then the matrix $e^{\underline{A}t}$ may contain numbers and functions which are complex in form and must be combined to form a solution which contains only real numbers and real functions. The task is tedious by hand, so the program combines complex functions into a real form convenient to the user.

The generation of a real form for $e^{\underline{A}t}$ with complex eigenvalues is handled in the same manner as in the case of a single linear equation with constant coefficients. Namely, if $\lambda_j = \alpha + i\beta$ is a characteristic root then so is $\lambda_\ell = \alpha - i\beta$, and the two roots have the same multiplicity. The term

$$\frac{t^{k-1}}{(k-1)!} \left(e^{\alpha t} e^{i\beta t} \underline{Z}_{j,k} + e^{\alpha t} e^{-i\beta t} \underline{Z}_{\ell,k} \right) \quad (48)$$

⁶Frame, loc. cit.

occurs in $e^{\frac{A}{t}t}$. Write

$$\underline{Z}_{j,k} = \underline{Z} + i\underline{ZZ} \quad (49)$$

$$\underline{Z}_{\ell,k} = \underline{W} + i\underline{WW} \quad (50)$$

with \underline{W} , \underline{WW} , \underline{Z} , \underline{ZZ} real $n \times n$ matrices, and write $e^{\pm i\beta t}$ as

$$\cos \beta t \pm i \sin \beta t \quad (51)$$

Multiplying out the complex numbers and grouping terms, Equation (48) becomes

$$\frac{t^{k-1} e^{\alpha t}}{(k-1)!} \left[\cos \beta t (\underline{Z} + \underline{W}) + \sin \beta t (\underline{WW} - \underline{ZZ}) + i \{ \sin \beta t (\underline{Z} - \underline{W}) + \cos \beta t (\underline{ZZ} + \underline{WW}) \} \right] \quad (52)$$

but $e^{\frac{A}{t}t}$ is real since \underline{A} is, therefore, the complex part of Equation (52) must vanish for all t . By evaluating Equation (52) at 0 and π/β , the connections

$$\underline{Z} = \underline{W}$$

$$\underline{ZZ} = -\underline{WW} \quad (53)$$

are established. Simplifying Equation (52) gives

$$\frac{t^{k-1}}{(k-1)!} e^{\alpha t} [2 \cos \beta t \cdot \underline{Z} - 2 \sin \beta t \cdot \underline{ZZ}] \quad (54)$$

All the quantities in Equation (54) are real.

It will now be shown how the complex computations of Equation (45) may be handled using only real computations. Recall that complex numbers can be changed into 2×2 matrices according to the following rule:

$$\alpha + i\beta \longmapsto \begin{pmatrix} \alpha & -\beta \\ \beta & \alpha \end{pmatrix} \quad (55)$$

If z and w are complex and \underline{Z} , \underline{W} are their corresponding matrices then^{7,8}

⁷Frame, op. cit.

⁸Herstein, loc. cit.

$$\begin{aligned}
z + w &\longmapsto \underline{z} + \underline{w} \\
zw &\longmapsto \underline{z}\underline{w} \\
z^{-1} &\longmapsto \underline{z}^{-1}
\end{aligned}
\tag{56}$$

Equations (56) are sufficient to guarantee that algebraic calculations done with the matrices will agree with those done with their complex numbers. This monomorphism is extended to $n \times n$ matrices in the natural way. If \underline{z} is a complex $n \times n$ matrix and

$$\underline{z} = \underline{A} + i\underline{B} \quad , \tag{57}$$

then

$$\underline{z} \longmapsto \begin{pmatrix} \underline{A} & -\underline{B} \\ \underline{B} & \underline{A} \end{pmatrix} \quad , \tag{58}$$

a $2n \times 2n$ matrix.

Similarly, if \underline{w} , \underline{c} , and \underline{d} are $n \times 1$ matrices such that

$$\underline{w} = \underline{c} + i\underline{d} \quad , \tag{59}$$

then

$$\underline{w} \longmapsto \begin{pmatrix} \underline{c} \\ \underline{d} \end{pmatrix} \quad . \tag{60}$$

Equations analogous to Equation (56) hold.^{9,10} Under these transformations, Equation (59) becomes

⁹Frame, loc. cit.

¹⁰Bradon, G. E., Introduction to Compact Transformation Groups, Academic Press, New York, New York, 1972.

$$\begin{pmatrix} \underline{I} \\ \underline{A} \\ \vdots \\ \underline{A}^{n-1} \\ 0 \\ \vdots \\ 0 \end{pmatrix} = \begin{pmatrix} \underline{VR}^t & -\underline{VC}^t \\ \underline{VC}^t & \underline{VR}^t \end{pmatrix} \begin{pmatrix} \underline{ZR}_{1,1} \\ \vdots \\ \underline{ZR}_{n,n_s} \\ \underline{ZI}_{1,1} \\ \vdots \\ \underline{ZI}_{n,n_s} \end{pmatrix} \quad (61)$$

where $\underline{V}^t = \underline{VR}^t + i\underline{VC}^t$ and $\underline{Z}_{i,j} = \underline{ZR}_{i,j} + i\underline{ZI}_{i,j}$. The $2n \times 2n$ matrix

$$\begin{pmatrix} \underline{VR}^t & -\underline{VC}^t \\ \underline{VC}^t & \underline{VR}^t \end{pmatrix} \quad (62)$$

is inverted and the calculations are made as before.

The $2n \times 2n$ matrix given in Equation (62) cannot be inverted by inverting the two $n \times n$ matrices \underline{VR} and \underline{VC} as can be shown by this counter-example:

$$\underline{A} = \begin{pmatrix} 0 & 1 & 0 \\ 0 & 0 & 1 \\ 0 & -1 & 0 \end{pmatrix} ; \quad (63)$$

hence

$$\text{Det}(\underline{A} - \lambda \underline{I}) = -(\lambda^3 + \lambda) \quad (64)$$

so

$$\lambda = 0, \pm i . \quad (65)$$

\underline{V}^t is

$$\begin{pmatrix} 1 & 1 & 1 \\ 0 & i & -i \\ 0 & -1 & -1 \end{pmatrix} . \quad (66)$$

The real matrix corresponding to \underline{V}^t is

$$\left(\begin{array}{ccc|ccc} 1 & 1 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & -1 & 1 \\ 0 & -1 & -1 & 0 & 0 & 0 \\ \hline 0 & 0 & 0 & 1 & 1 & 1 \\ 0 & 1 & -1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & -1 & -1 \end{array} \right) \quad (67)$$

None of the 3×3 matrices are invertible.

e. Error Checks

Recall that $e^{\underline{A}t}$ satisfies

$$\frac{d e^{\underline{A}t}}{dt} = \underline{A} e^{\underline{A}t} \quad (68)$$

Furthermore, among all the analytic functions satisfying Equation (68) it alone satisfies

$$f(0) = \underline{I} \quad (69)$$

Equations (68) and (69) can be used to construct a test on the validity of any technique purporting to calculate $e^{\underline{A}t}$ in terms of its constituent matrices. If $e^{\underline{A}t}$ is differentiated with respect to t [Equation (15)], Equation (70) results

$$\sum_{j=1}^s \sum_{k=1}^{n_j} \frac{t^{k-1} e^{\lambda_j t}}{(k-1)!} [Z_{j,k+1} + \lambda_j Z_{j,k}] \quad (70)$$

(without loss of generality we can set $Z_{j,n_j+1} = 0$). In terms of Equation (68)

$$\sum_{j=1}^s \sum_{k=1}^{n_j} \frac{t^{k-1} e^{\lambda_j t}}{(k-1)!} [Z_{j,k+1} + \lambda_j Z_{j,k} - \underline{A} Z_{j,k}] = 0 \quad (71)$$

or setting

$$\frac{1}{(k-1)!} Z_{j,k} = Z'_{j,k} \quad (72)$$

$$\sum_{j=1}^s \sum_{k=1}^{n_j} t^{k-1} e^{\lambda_j t} \left[k \underline{Z}_{j,k+1} + \lambda_j \underline{Z}_{j,k} - \underline{AZ}_{j,k} \right] = \underline{0} \quad (73)$$

The functions $t^{k-1} e^{\lambda_j t}$ are linearly independent so their coefficient must be zero for the left hand side of Equation (73) to equal the null matrix:

$$\begin{aligned} k \underline{Z}'_{j,k+1} + \lambda_j \underline{Z}'_{j,k} - \underline{AZ}'_{j,k} &= \underline{0} \quad k < n_j \\ \lambda_j \underline{Z}'_{j,n_j} - \underline{AZ}'_{j,n_j} &= \underline{0} \end{aligned} \quad (74)$$

If $\lambda_j = \alpha + i\beta$ is complex then, using analogous notation,

$$\begin{aligned} k \underline{ZR}'_{j,k+1} + \alpha \underline{ZR}'_{j,k} + \beta \underline{ZI}'_{j,k} - \underline{AZR}'_{j,k} &= \underline{0} \\ k \underline{ZI}'_{j,k+1} + \alpha \underline{ZI}'_{j,k} - \beta \underline{ZR}'_{j,k} - \underline{AZI}'_{j,k} &= \underline{0} \\ \alpha \underline{ZR}'_{j,n_j} + \beta \underline{ZI}'_{j,n_j} - \underline{AZR}'_{j,n_j} &= \underline{0} \\ \alpha \underline{ZI}'_{j,n_j} - \beta \underline{ZR}'_{j,n_j} - \underline{AZI}'_{j,n_j} &= \underline{0} \end{aligned} \quad (75)$$

Checking these equations provides a good overall check for numerical errors. The program keeps track of the maximum entry of the matrix on the left hand sides of Equations (74) and (75) (absolute value of the matrix entry). After all calculations have been done, it prints out this maximum.

3. Program Description

a. Flow Chart of Program

A flow chart of the program containing the computations defined in Section 2 is given in Figure 1. All blocks in the flow chart with the exception of the last, are either self-explanatory or have already been explained. In calculating the inverse of a matrix the IBM subroutine INVERT,¹¹ which employs pivotal condensation, was used. In

¹¹System/360 Scientific Subroutine Package, International Business Machines, White Plains, New York, 1968.

calculating the roots of the characteristic polynomial POLRT, another IBM routine was used.¹² It is a 500-step Newton-Raphson method in two variables.

b. Inputs to Program

This section defines the form of the input data for the program. The first four cards define the tolerance constants described in Section 2. These cards are read only once for any set of A matrices to be run. If the dimension card (LDEM) for an A matrix is zero in value, the computer run will terminate.

The order and format of the input data cards are shown in Figure 2. The variable names in Figure 2 can be identified as follows:

DB - In sorting roots, two roots which differ by an amount less than DB are set equal. This tolerance will effect the multiplicity calculated for a root and hence the form of the solution. If the user selects DB = 0, distinct roots are likely to be generated changing the form of the solution. Since the solution to Equation (1) depends continuously on the eigenvalues, the form of the solution chosen by the program will give the correct answer in a numeric sense.

DD - If the determinant of the Vandermonde is less than DD in absolute value, an error message is printed stating that the Vandermonde matrix is singular. When this message is encountered, one of the following occurred: (1) DD was selected too large, (2) User entered data improperly, or (3) numerical roundoff or truncation error is the source of the problem. After the error message is printed, execution is halted and the next data set is read in (starting with LDEM). If the user sets DD = 0, the program halts only if the determinant of V is zero.

LDEM - Order of the system matrix.

A(I,J) - Element of the system matrix.

c. Sample Program Output

Along with the error messages already discussed, the program will print the A matrix, its characteristic polynomial, a table of characteristic roots and their multiplicities, $e^{\underline{A}t}$ expressed as a sum of constituent matrices and analytic functions, and the maximum error as determined by the method of Section 2.e.

¹²Ibid.

The input data set shown in Figure 3 was used to generate the output shown in Figure 4.

4. Program Listing

The listing given in Figure 5 is the double precision version of the EAT program.

5. Conclusions

Programs to calculate the state transition matrix (e^{At}) as a linear combination of functions of time are not generally available. The program described in the report computes e^{At} and should be useful to those involved in the solution of linear differential equations described by state space equations.

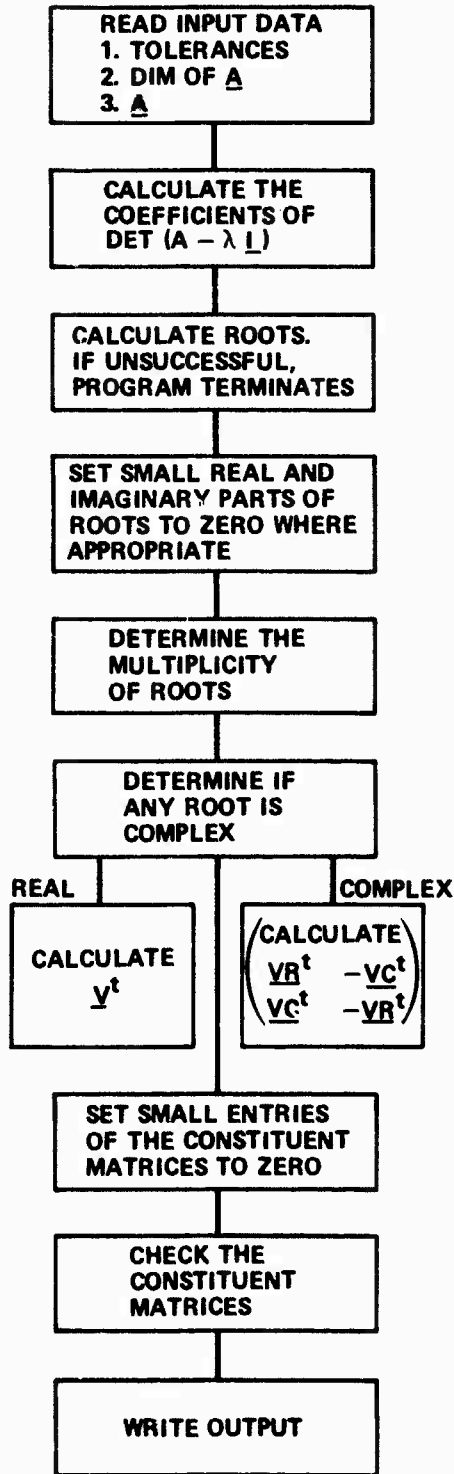


Figure 1. Flow chart of program.

	<u>CARD NO.</u>	<u>VARIABLE NAME</u>	<u>FORMAT</u>
TOLERANCE LEVELS	{ 1	DB	D15.0
	2	DC	D15.0
	3	DD	D15.0
FIRST DATA SET	{ 4	LDEM	I2
	5	A(1,1), A(1,2), A(1,3), A(1,4)	4D15.0
	.	.	.
	.	.	.
SECOND DATA SET	{ M	LDEM	I2
	M+1	A(1,1), A(1,2), A(1,3), A(1,4)	4D15.0
	.	.	.
	.	.	.
RUN TERMINATION	{ K	BLANK CARD	I2

Figure 2. Input data card format.

5	15	30	45	60	← CARD COLUMN NUMBER
0.0001					
0.0001					
0.0001					
0.0001					
4					
0.0	1.0	0.0	0.0	}	1ST DATA SET
0.0	0.0	1.0	0.0		
0.0	0.0	0.0	1.0		
-1.0	0.0	-2.0	0.0		
3					
0.0	1.0	0.0	0.0	}	2ND DATA SET
0.0	1.0	-4.0	-6.0		
-4.0					
4					
0.0	1.0	0.0	0.0	}	3RD DATA SET
0.0	0.0	1.0	0.0		
0.0	0.0	0.0	1.0		
-4.0	-8.0	-8.0	-4.0		
4					
0.0	0.0	1.0	0.0	}	4TH DATA SET
1.0	0.0	0.0	1.0		
0.0	0.0	0.0	-1.0		
0.0	1.0	0.0	0.0		
(BLANK CARD)					

Figure 3. Sample input data card set.

1st Data Set Output

A MATRIX

0.	.100000000D+01	0.	0.
0.	0.	.100000000D+01	0.
0.	0.	0.	.100000000D+01
-.100000000D+01	0.	-.200000000D+01	0.

CHARACTERISTIC POLYNOMIAL

(.10000000D+01) +
 (-0.) *X** 1 +
 (.20000000D+01) *X** 2 +
 (-0.) *X** 3 +
 (-.10000000D+01) *X** 4

CHARACTERISTIC ROOTS

REAL PART	COMPLEX PART	MULTIPLICITY
0.	-.10000000D+01	2
0.	.10000000D+01	2

EXP(AX) =

+ F(1)*Z(1, 1) + F(2)*ZZ(1, 1)
 + F(3)*Z(1, 2) + F(4)*ZZ(1, 2)

WHERE F(I) ARE

Figure 4. Sample program output.

F(1) = COS(X* -.10000000D+01)

F(2) = SIN(X* -.10000000D+01)

F(3) = (X** 1)*COS(X* -.10000000D+01)

F(4) = (X** 1)*SIN(X* -.10000000D+01)

AND WHERE THE Z AND ZZ MATRICES ARE

Z(1, 1) MATRIX IS

.100000000D+01	0.	0.	0.
0.	.100000000D+01	0.	0.
0.	0.	.100000000D+01	0.
0.	0.	0.	.100000000D+01

ZZ(1, 1) MATRIX IS

0.	-.150000000D+01	0.	-.500000000D+00
.500000000D+00	0.	-.500000000D+00	0.
0.	.500000000D+00	0.	-.500000000D+00
.500000000D+00	0.	.150000000D+01	0.

Z(1, 2) MATRIX IS

0.	-.500000000D+00	0.	-.500000000D+00
.500000000D+00	0.	.500000000D+00	0.
0.	.500000000D+00	0.	.500000000D+00
-.500000000D+00	0.	-.500000000D+00	0.

ZZ(1, 2) MATRIX IS

-.500000000D+00	0.	-.500000000D+00	0.
0.	-.500000000D+00	0.	-.500000000D+00
.500000000D+00	0.	.500000000D+00	0.
0.	.500000000D+00	0.	.500000000D+00

MAXIMUM ENTRY OF ERROR MATRIX= .40766752D-12

.....

Figure 4. (Continued).

2nd Data Set Output

A MATRIX

0.	.100000000D+01	0.
0.	0.	.100000000D+01
-.400000000D+01	-.600000000D+01	-.400000000D+01

CHARACTERISTIC POLYNOMIAL

(.40000000D+01) +
 (.60000000D+01) *X** 1 +
 (.40000000D+01) *X** 2 +
 (.10000000D+01) *X** 3

CHARACTERISTIC ROOTS

REAL PART	COMPLEX PART	MULTIPLICITY
-.20000000D+01	0.	1
-.10000000D+01	-.10000000D+01	1
-.10000000D+01	.10000000D+01	1

EXP(AX) =

+ F(1)*Z(1, 1)
 + F(2)*Z(2, 1) + F(3)*Z(2, 1)

WHERE F(I) ARE

F(1) = EXP(X* -.20000000D+01) -

Figure 4. (Continued).

$$F(2) = (X**0)*EXP(X*-.10000000D+01)*COS(X*-.10000000D+01)$$

$$F(3) = (X**0)*EXP(X*-.10000000D+01)*SIN(X*-.10000000D+01)$$

AND WHERE THE Z AND ZZ MATRICES ARE

Z(1,1) MATRIX IS

.100000000D+01	.100000000D+01	.500000000D+00
-.200000000D+01	-.200000000D+01	-.100000000D+01
.400000000D+01	.400000000D+01	.200000000D+01

Z(2,1) MATRIX IS

0.	-.100000000D+01	-.500000000D+00
.200000000D+01	.300000000D+01	.100000000D+01
-.400000000D+01	-.400000000D+01	-.100000000D+01

ZZ(2,1) MATRIX IS

-.200000000D+01	-.200000000D+01	-.500000000D+00
.200000000D+01	.100000000D+01	0.
0.	.200000000D+01	.100000000D+01

MAXIMUM ENTRY OF ERROR MATRIX= .60584518D-27

Figure 4. (Continued).

3rd Data Set Output

A MATRIX

0.	.100000000D+01	0.	0.
0.	0.	.100000000D+01	0.
0.	0.	0.	.100000000D+01
-.400000000D+01	-.800000000D+01	-.800000000D+01	-.400000000D+01

CHARACTERISTIC POLYNOMIAL

(.400000000D+01) *
 (.800000000D+01) *X** 1 *
 (.800000000D+01) *X** 2 *
 (.400000000D+01) *X** 3 *
 (.100000000D+01) *X** 4

CHARACTERISTIC ROOTS

REAL PART	COMPLEX PART	MULTIPLICITY
-.100000000D+01	-.100000000D+01	2
-.100000000D+01	.100000000D+01	2

EXP(AX) =

+ F(1)*Z(1, 1) + F(2)*ZZ(1, 1)
 + F(3)*Z(1, 2) + F(4)*ZZ(1, 2)

WHERE F(I) ARE

Figure 4. (Continued).

$$F(1) = (X**0)*EXP(X*-.10000000D+01)*COS(X*-.10000000D+01)$$

$$F(2) = (X**0)*EXP(X*-.10000000D+01)*SIN(X*-.10000000D+01)$$

$$F(3) = (X**1)*EXP(X*-.10000000D+01)*COS(X*-.10000000D+01)$$

$$F(4) = (X**1)*EXP(X*-.10000000D+01)*SIN(X*-.10000000D+01)$$

AND WHERE THE Z AND ZZ MATRICES ARE

Z(1,1) MATRIX IS

.100000000D+01	0.	0.	0.
0.	.100000000D+01	0.	0.
0.	0.	.100000000D+01	0.
0.	0.	0.	.100000000D+01

ZZ(1,1) MATRIX IS

-.200000000D+01	-.300000000D+01	-.150000000D+01	-.500000000D+00
.200000000D+01	.200000000D+01	.100000000D+01	.500000000D+00
-.200000000D+01	-.200000000D+01	-.200000000D+01	-.100000000D+01
.400000000D+01	.600000000D+01	.600000000D+01	.200000000D+01

Z(1,2) MATRIX IS

-.100000000D+01	-.200000000D+01	-.150000000D+01	-.500000000D+00
.200000000D+01	.300000000D+01	.200000000D+01	.500000000D+00
-.200000000D+01	-.200000000D+01	-.100000000D+01	0.
0.	-.200000000D+01	-.200000000D+01	-.100000000D+01

ZZ(1,2) MATRIX IS

-.100000000D+01	-.100000000D+01	-.500000000D+00	0.
0.	.100000000D+01	.100000000D+01	-.500000000D+00
.200000000D+01	.400000000D+01	.300000000D+01	.100000000D+01
-.400000000D+01	-.600000000D+01	-.400000000D+01	-.100000000D+01

MAXIMUM ENTRY OF ERROR MATRIX= .10020293D-10

Figure 4. (Continued).

4th Data Set Output

A MATRIX

0.	0.	.100000000D+01	0.
.100000000D+01	0.	0.	.100000000D+01
0.	0.	0.	-.100000000D+01
0.	.100000000D+01	0.	0.

CHARACTERISTIC POLYNOMIAL

(.100000000D+01) +

(-0.) *X** 1 +

(-.100000000D+01) *X** 2 +

(-0.) *X** 3 +

(.100000000D+01) *X** 4

CHARACTERISTIC ROOTS

REAL PART	COMPLEX PART	MULTIPLICITY
-.86602540D+00	.500000000D+00	1
-.86602540D+00	-.500000000D+00	1
.86602540D+00	-.500000000D+00	1
.86602540D+00	.500000000D+00	1

EXP(AX) =

+ F(1)*Z(1, 1) + F(2)*ZZ(1, 1)

+ F(3)*Z(3, 1) + F(4)*ZZ(3, 1)

Figure 4. (Continued).

WHERE F(I) ARE

$$F(1) = (X**0)*EXP(X*-.866025400+00)*COS(X*.500000000+00)$$

$$F(2) = (X**0)*EXP(X*-.866025400+00)*SIN(X*.500000000+00)$$

$$F(3) = (X**0)*EXP(X*.866025400+00)*COS(X*-.500000000+00)$$

$$F(4) = (X**0)*EXP(X*.866025400+00)*SIN(X*-.500000000+00)$$

AND WHERE THE Z AND ZZ MATRICES ARE

Z(1,1) MATRIX IS

.500000000+00	-.2886751350+00	-.5773502690+00	0.
-.2886751350+00	.5000000000+00	0.	-.5773502690+00
-.2886751350+00	0.	.5000000000+00	.2886751350+00
0.	-.2886751350+00	.2886751350+00	.5000000000+00

ZZ(1,1) MATRIX IS

.2886751350+00	-.5000000000+00	0.	.5773502690+00
.5000000000+00	-.2886751350+00	-.5773502690+00	0.
-.5000000000+00	.5773502690+00	.2886751350+00	-.5000000000+00
-.5773502690+00	.5000000000+00	.5000000000+00	-.2886751350+00

Z(3,1) MATRIX IS

.5000000000+00	.2886751350+00	.5773502690+00	0.
.2886751350+00	.5000000000+00	0.	.5773502690+00
.2886751350+00	0.	.5000000000+00	-.2886751350+00
0.	.2886751350+00	-.2886751350+00	.5000000000+00

ZZ(3,1) MATRIX IS

.2886751350+00	.5000000000+00	0.	.5773502690+00
-.5000000000+00	-.2886751350+00	-.5773502690+00	0.
.5000000000+00	.5773502690+00	.2886751350+00	.5000000000+00
-.5773502690+00	-.5000000000+00	-.5000000000+00	-.2886751350+00

MAXIMUM ENTRY OF ERROR MATRIX= .378653230-28

Figure 4. (Concluded).

```

PROGRAM EAT(INPUT,OUTPUT,TAPE5=INPUT,TAPE6=OUTPUT)
DIMENSION XCOF(10),COF(10),ROOTR(10),ROOTI(10)
DIMENSION STORX(10),STORY(10)
DIMENSION VR(10,10),VC(10,10),RSAV(10,10)
5 DIMENSION KF(10),N(10),RELL(10),CMPXX(10)
DIMENSION A(10,10),Z(10,10),ZZ(10,10),PROD(20,20)
DIMENSION V(20,20),ST400
DIMENSION L1(20),M1(20)
DOUBLE PRECISION RELL,CMPXX,A,Z,ZZ,PROD
10 DDUBLE PRECISION V,S,O,B
DDUBLE PRECISION DA,DB,DC,DD
DDUBLE PRECISION XCOF,COF,ROOTR,ROOTI,STORY,STORX
DDUBLE PRECISION Y,X
DDUBLE PRECISION YY
15 DDUBLE PRECISION VR,VC,RSAV
C
C READ IN ALL THE DATA
C
20 READ(5,101) DB
READ(5,101) DC
READ(5,101) DD
101 FORMAT(10I5.0)
866 WRITE(6,755)
WRITE(6,937)
25 937 FORMAT(60(IK,1H*))
WRITE(6,801)
801 FORMAT(1HI)
READ(5,I) LDEM
K=0
30 IFF=0
IF(LDEM) 761,762,761
761 READ(5,88)((A(I,J),J=1,LDEM),I=1,LDEM)
88 FORMAT(40I5.0)
WRITE(6,60)
35 60 FORMAT(/////3X,BHA MATRIX)
CALL WRITE1(LDEM,A)
WRITE(6,755)
WRITE(6,937)
C
C CALCULATE THE CHARACTERISTIC POLYNOMIAL
C
XCOF(LDEM+1)=1.00
CALL TRACE(A,LDEM,B)
XCOF(LDEM)=-B
45 DO 40 I=1,LDEM
DO 40 J=1,LDEM
40 VR(I,J)=A(I,J)
DO 10 I=1,LDEM
10 VR(I,I)=VR(I,I)+XCOF(LDEM)
50 MM=LDEM-1
M=LDEM
DO 33 I=1,MM
CALL MATPRO(M,A,VR,PROD)
CALL TRACE2(PROD,M,B)
55 XCOF(M-1)=-((1.00/DBLE(1.00+1.00)))*B
DO 20 IP=1,M
DO 20 JP=1,M

```

Figure 5. Program listing.

```

PROGRAM EAT          74/74  OPT=I          FTN 4.2+75067      05/16/75  00.59.29.

20 VR(IP,JP)=PRODTIP,JP)
DO 30 IP=1,M
60 30 VR(IP,IP)=VR(IP,IP)+XCOF(M-I)
33 CONTINUE
WRITE(6,755)
WRITE(6,104)
104 FORMAT(3X,25HCHARACTERISTIC POLYNOMIAL//)
65 CALL WRITEP(XCOF,LOEM)
C
C CALCULATE THE CHARACTERISTIC ROOTS
C
CALL POLRT(XCOF,COF,LOEM,ROOTR,ROOTI,IER)
70 DO 399 I=1,LOEM
IF (ROOTR(I).LT.DB.AND.ROOTR(I).GT.-DB) ROOTR(I)=0.00
IF (ROOTI(I).LT.DB.AND.ROOTI(I).GT.-DB) ROOTI(I)=0.00
399 CONTINUE
IF (IER) 61,70,61
75 61 WRITE(6,108) IER
108 FORMAT(///3X,30HERROR IN ROOT CALCULATION,MODE,12)
GO TO 866
C
C SORT AND CLASSIFY THE ROOTS
C
80 70 CALL SORT(ROOTR,ROOTI,RELL,CMPXX,STORX,STORY,N,LOEM,K,DB)
WRITE(6,755)
755 FORMAT(3X,///)
WRITE(6,55)
85 55 FORMAT(2X,20HCHARACTERISTIC ROOTS //2X,9HREAL PART,13X,12HCOMPLEX
1 PART,13X,12HMULTIPLICITY )
IQ=K
DO 700 I=1,IQ
700 WRITE(6,701) RELL(I),CMPXX(I),N(I)
90 701 FORMAT(1X,D15.8,5X,D15.8,15X,15)
WRITE(6,755)
WRITE(6,937)
C
C CALCULATE AND INVERT THE VANDERMONDE
C
95 KF(I)=I
DO 883 I=1,LOEM
KFX=I+I
883 KF(I+I)=KFX*KF(I)
DO 313 III=1,K
100 IF (CMPXX(III).NE.0.00) GO TO 129
313 CONTINUE
CALL VANDR(RELL,LOEM,N,K,V)
GO TO 107
105 129 LOEM2=2*LOEM
CALL VAND(RELL,CMPXX,LOEM,N,K,V)
CALL ARRAY(2,LOEM2,LOEM2,20,20,S,V)
CALL GINVRT(S,LOEM2,0,LI,MI)
CALL ARRAY(1,LOEM2,LOEM2,20,20,S,V)
110 D=DABS(D)
O=OSORT(D)
IF (D.GT.00.OR.D.LT.-00) GO TO 846
647 WRITE(6,648)
648 FORMAT(3X,40HV MATRIX IS SINGULAR -- CHECK INPUT DATA)

```

Figure 5. (Continued).

```

115      GO TO 866
      646 KFG=1
      GO TO 555
107      CALL ARRAY(2,LOEM,LOEM,20,20,5,V)
      CALL GINVRT(5,LOEM,D,LI,M1)
120      CALL ARRAY(1,LOEM,LOEM,20,20,5,V)
      IF(D.LT.DD.AND.D.GT.-DD) GO TO 847
      C
      C      DETERMINE THE FORM OF EXP(AT)
      C
125      649 KFG=0
      LOEM2=LOEM
      555 IK=0
      WRITE(6,756)
      756 FORMAT(3X,///)
130      WRITE(6,411)
      IIK=0
      KOOP=0
      DO 444 I=1,K
      IF(KOOP) 414,415,414
135      414 KOOP=0
      GO TO 444
      415 LGM=N(I)
      DO 400 J=1,LGM
      IIK=IIK+1
140      IF(CMPAX(I)) 406,401,406
      401 WRITE(6,410) IIK,I,J
      GO TO 400
      406 LOG=IIK+1
      WRITE(6,440) IIK,I,J,LOG,I,J
145      IIK=LOG
      KOOP=1
      400 CONTINUE
      444 CONTINUE
      411 FORMAT(3X,94EXP(AX) = /)
150      440 FORMAT(3X,4H* F(,I2,4H)*Z(,I2,1H,,I2,6H) + F(,I2,5H)*ZZ(,I2,1H,,
      I2,1H)/)
      410 FORMAT(3X,4H* F(,I2,4H)*Z(,I2,1H,,I2,1H)/)
      WRITE(6,430)
      KOOP=0
      IIK=0
155      DO 901 I=1,K
      LGM=N(I)
      IF(KOOP) 614,615,614
      614 KOOP=0
      GO TO 901
160      615 DO 600 J=1,LGM
      IIK=IIK+1
      IF(CMPAX(I)) 606,601,606
      601 IF(J=1) 604,602,604
165      602 WRITE(6,603) IIK,RELL(I)
      GO TO 600
      604 JJJ=J-1
      WRITE(6,605) IIK,JJJ,RELL(I)
      GO TO 600
170      606 KOOP=1
      IF(RELL(I)) 612,607,612

```

Figure 5. (Continued).

```

607 IF(J=1) 610,608,610
608 WRITE(6,609) I1K,CMPXX(1)
      I1K=I1K+1
175   WRITE(6,709) I1K,CMPXX(1)
      KOOP=1
      GO TO 600
610 JJJ=J-1
      WRITE(6,611) I1K,JJJ,CMPXX(1)
      I1K=I1K+1
180   WRITE(6,711) I1K,JJJ,CMPXX(1)
      GO TO 600
612 JJJ=J-1
      WRITE(6,613) I1K,JJJ,RELL(I),CMPXX(1)
      I1K=I1K+1
185   WRITE(6,713) I1K,JJJ,RELL(1),CMPXX(1)
600 CONTINUE
901 CONTINUE
605 FORMAT(///3X,2HF(,12.8H) = (X*,12.8H)*EXP(X*,015.8,1H))
190 611 FORMAT(///3X,2HF(,12.8H) = (X*,12.8H)*COS(X*,015.8,1H))
611 FORMAT(///3X,2HF(,12.8H) = (X*,12.8H)*SIN(X*,015.8,1H))
613 FORMAT(///3X,2HF(,12.8H) = (X*,12.8H)*EXP(X*,015.8,
18H)*COS(X*,015.8,1H))
195 613 FORMAT(///3X,2HF(,12.8H) = (X*,12.8H)*EXP(X*,015.8,
18H)*SIN(X*,015.8,1H))
603 FORMAT(///3X,2HF(,12.10H) = EXP(X*,015.8,1H))
609 FORMAT(///3X,2HF(,12.4H) = ,6HCOS(X*,015.8,1H))
709 FORMAT(///3X,2HF(,12.4H) = ,6HSIN(X*,015.8,1H))
200 430 FORMAT(///3X,14HWHERE F(1) ARE)
      WRITE(6,760)
760 600 FORMAT(///3X,35HANO WHERE THE Z AND ZZ MATRICES ARE )
      C
      C
      C
205   KOOP=0
      IK=0
      YY=0.00
      DO 809 I=1,K
      IF(KOOP) 616,617,616
210 616 KOOP=0
      IK=IK+1
      GO TO 809
617 MMM=N(I)
      DO 109 J=1,MMM
215   IK=IK+1
      IF(CMPXX(1)) 618,619,618
618 KOOP=1
      KFG=1
      GO TO 620
220 619 KFG=0
620 DO 111 L1L=1,LDEM
      DO 111 L2L=1,LDEM
      IF(L1L=L2L) 113,112,113
112 Z(L1L,L2L)=V(IK,1)
      GO TO 111
225 113 Z(L1L,L2L)=0.0
111 CONTINUE
      IF(KFG) 205,206,205

```

Figure 5. (Continued).

```

230      205 DO 114 ILL=1,LDEN
          DO 114 JLL=1,LDEN
          IF (ILL-JLL) 116,115,116
          115 MKM=IK+LDEN
          ZZ(ILL,JLL)=V(MKM,1)
          GO TO 114
235      116 ZZ(ILL,JLL) =0.0
          114 CONTINUE
          206 DO 140 JL=2,LDEN
          IF (JL.EQ.2) GO TO 173
          IF (JL.EQ.3) GO TO 132
240      CALL MATPRO(LDEN,A,RSAB,PROD)
          GO TO 273
          132 DO 236 IP=1,LDEN
          DO 236 JP=1,LDEN
          236 RSAB(IP,JP)=A(IP,JP)
245      CALL MATPRO(LDEN,A,RSAB,PROD)
          273 DO 199 IP=1,LDEN
          DO 199 JP=1,LDEN
          199 RSAB(IP,JP)=PRDD(IP,JP)
          GO TO 120
250      173 DO 919 IP=1,LDEN
          DO 919 JP=1,LDEN
          919 PROD(IP,JP)=A(IP,JP)
          120 DO 119 IN=1,LDEN
          DO 119 JN=1,LDEN
255      Z(IN,JN)=Z(IN,JN)+V(IK,JL)*PRDD(IN,JN)
          IF (KFG) 333,119,333
          333 LKL=IK+LDEN
          ZZ(IN,JN)=ZZ(IN,JN)+V(LKL,JL)*PROD(IN,JN)
          119 CONTINUE
260      140 CONTINUE
          IF (KFG) 758,757,758
          758 DO 759 MIN=1,LDEN
          DO 759 MJN=1,LDEN
          Z(MIN,MJN)=2.00*Z(MIN,MJN)
265      759 ZZ(MIN,MJN)=2.00*ZZ(MIN,MJN)
          757 WRITE(6,776) I,J
          IF (J.EQ.1) X=1.00
          IF (J.GT.1) X=DBLE(FLOAT(KF(J-1)))
          DO 827 MU=1,LDEN
          DO 827 MUT=1,LDEN
270      Z(MU,MUT)=Z(MU,MUT)/X
          827 CONTINUE
          DO 321 MU=1,LDEN
          DO 321 MUT=1,LDEN
275      X=Z(MU,MUT)
          IF (X.LT.DC.AND.X.GT.-DC) Z(MU,MUT)=0.00
          321 CONTINUE
          CALL WHITE1 (LDEN,Z)
          IF (KFG) 207,309,207
280      X=RELL(I)
          IF (J.EQ.1) GO TO 157
          CALL MATPRO(LDEN,A,VR,PROD)
          DO 837 IP=1,LDEN
          DO 837 JP=1,LDEN
285      Y=X*VR(IP,JP)+Z(IP,JP)*FLOAT(J-1)-PROD(IP,JP)

```

Figure 5. (Continued).

```

      Y=OABS(Y)
      14 IF(Y-YY) 837,17,17
      17 YY=Y
      837 CONTINUE
290 157 DO 783 IP=1,LOEM
      DO 783 JP=1,LOEM
      783 VR(IP,JP)=Z(IP,JP)
      IF(J.NE.MMM) GO TO 109
      CALL MATPRO(LOEM,A,VR,PROO)
      295 DO 387 IP=1,LOEM
      DO 387 JP=1,LOEM
      Y=PROO(IP,JP)-VR(IP,JP)*X
      Y=OABS(Y)
      15 IF(Y-YY) 387,16,16
      16 YY=Y
      387 CONTINUE
      856 IF(KFG) 207,109,207
      207 WRITE(6,779) I,J
      IF(J.EQ.I) X=I.00
      305 IF(J.GT.I) X=OBLE(FLOAT(KF(J-1)))
      DO 839 MU=1,LOEM
      DO 839 MUT=1,LOEM
      ZZ(MU,MUT)=ZZ(MU,MUT)/X
      839 CONTINUE
      310 DO 128 MU=1,LOEM
      DO 128 MUT=1,LOEM
      X=ZZ(MU,MUT)
      IF(X.LT.DC.AND.X.GT.-DC) ZZ(MU,MUT)=0.00
      128 CONTINUE
      315 CALL WRITE1(LOEM,ZZ)
      IF(J.EQ.I) GO TO 717
      CALL MATPRO(LOEM,A,VR,PROO)
      DO 703 IP=1,LOEM
      DO 703 JP=1,LOEM
      320 Y=FLOAT(J-1)*Z(IP,JP)+RELL(I)*VR(IP,JP)+CMPXX(I)*VC(IP,JP)
      Y=Y-PROO(IP,JP)
      Y=OABS(Y)
      21 IF(Y-YY) 703,22,22
      22 YY=Y
      325 703 CONTINUE
      CALL MATPRO(LOEM,A,VC,PROO)
      DO 370 IP=1,LOEM
      DO 370 JP=1,LOEM
      Y=FLOAT(J-1)*ZZ(IP,JP)+RELL(I)*VC(IP,JP)+CMPXX(I)*VR(IP,JP)
      330 Y=Y-PROO(IP,JP)
      Y=OABS(Y)
      23 IF(Y-YY) 370,24,24
      24 YY=Y
      370 CONTINUE
      335 717 DO 307 IP=1,LOEM
      DO 307 JP=1,LOEM
      VR(IP,JP)=Z(IP,JP)
      VC(IP,JP)=ZZ(IP,JP)
      307 CONTINUE
      340 IF(J.NE.MMM) GO TO 109
      CALL MATPRO(LOEM,A,VR,PROO)
      DO 893 IP=1,LOEM
      DO 893 JP=1,LOEM
      Y=RELL(I)*VR(IP,JP)+CMPXX(I)*VC(IP,JP)+PROO(IP,JP)
      345 Y=OABS(Y)
      41 IF(Y-YY) 893,42,42
      42 YY=Y
      893 CONTINUE
      350 CALL MATPRO(LOEM,A,VC,PROO)
      DO 993 IP=1,LOEM
      DO 993 JP=1,LOEM
      Y=CMPXX(I)*VR(IP,JP)+RELL(I)*VC(IP,JP)+PROO(IP,JP)
      Y=OABS(Y)
      355 43 IF(Y-YY) 993,44,44
      44 YY=Y
      993 CONTINUE
      104 CONTINUE
      889 CONTINUE
      WRITE(6,755)
      360 WRITE(6,19) YY
      19 FORMAT(3X,30HMAXIMUM ENTRY OF ERROR MATRIX=,1015.8)
      GO TO 866
      179 FORMAT(///3X,3HZ(,I2,14,.,I2,11H) MATRIX IS/)
      778 FORMAT(///3X,2HZ(,I2,14,.,I2,11H) MATRIX IS/)
      365 1 FORMAT(I2)
      762 CONTINUE
      END

```

Figure 5. (Continued).

```

SUBROUTINE GINVRT(A,N,D,L,M)
C
C THIS ROUTINE CALCULATES THE INVERSE OF A MATRIX
C
5  DIMENSION A(I),L(I),M(I)
DOUBLE PRECISION A,BIGA,D,HOLD
D=1.
NK=-N
DO 80 K=1,N
10  NK=NK+N
    L(K)=K
    M(K)=K
    KK=NK+K
    BIGA=A(KK)
15  DO 20 J=K,N
    IZ=N*(J-I)
    DO 20 I=K,N
    IJ=IZ+I
20  IF (DABS(BIGA)-DABS(A(IJ))) 15,20,20
    BIGA=A(IJ)
    L(K)=I
    M(K)=J
    CONTINUE
20  C INTERCHANGE ROWS
    J=L(K)
25  IF (J=K) 35,35,25
    KI=K-N
    DO 30 I=1,N
    KI=KI+N
    HOLD=-A(KI)
30  JI=KI-K+J
    A(KI)=A(JI)
    A(JI)=HOLD
15  C INTERCHANGE COLUMNS
    I=M(K)
35  IF (I=K) 45,45,38
    JP=N*(I-I)
38  DO 40 J=1,N
    JK=NK+J
    JI=JP+J
40  HOLD=-A(JK)
    A(JK)=A(JI)
    A(JI)=HOLD
40  C DIVIDE COLUMN BY MINUS PIVOT (VALUE OF PIVOT ELEMENT IS CONTAINED
    C IN BIGA)
45  IF (BIGA) 49,46,48
46  O=0.
    RETURN
48  DO 55 I=1,N
    IF (I=K) 50,55,50
50  IK=NK+I
    A(IK)=A(IK)/(-BIGA)
55  CONTINUE
C REDUCE MATRIX
55  DO 65 I=1,N
    IK=NK+I
    HOLD=A(IK)

```

Figure 5. (Continued).


```

      IJ=I-N
      DO 65 J=1,N
      IJ=IJ+N
      IF(I-K) 60,65,60
      IF(J-K) 62,65,62
      KJ=IJ-I+K
      A(IJ)=MOLO*A(KJ)+A(IJ)
      CONTINUE
      C
      DIVIDE ROW BY PIVOT
      KJ=K-N
      DO 75 J=1,N
      KJ=KJ+N
      IF(J-K) 70,75,70
      A(KJ)=A(KJ)/BIGA
      CONTINUE
      C
      PRODUCT OF PIVOTS
      D=D*BIGA
      C
      REPLACE PIVOT BY RECIPROCAL
      A(KK)=1./BIGA
      CONTINUE
      C
      FINAL ROW AND COLUMN INTERCHANGE
      K=N
      K=(K-1)
      IF(K) 150,150,105
      I=L(K)
      IF(I-K) 120,120,108
      JQ=N*(K-1)
      JR=N*(I-1)
      DO 110 J=1,N
      JK=JQ+J
      MOLO=A(JK)
      JI=JR+J
      A(JK)=-A(JI)
      A(JI)=MOLO
      J=M(K)
      IF(J-K) 100,100,125
      KI=K-N
      DO 130 I=1,N
      KI=KI+N
      MOLO=A(KI)
      JI=KI-K+J
      A(KI)=-A(JI)
      A(JI)=MOLO
      GO TO 100
      RETURN
      END

```

Figure 5. (Continued).

SUBROUTINE ARRAY 74/74 OPT=1

FTN 4.2+75067 05/16/75 09.59.58.

```

      SUBROUTINE ARRAY (MODE,I,J,M,N,S,D)
      C
      C THIS ROUTINE PREPARES A MATRIX FOR SINVERT
      C
5     DIMENSION S(I,J,D(I))
      DOUBLE PRECISION S,D
      NI=N-1
      C TEST TYPE OF CONVERSION
      IF (MODE=1) 100,100,120
      C CONVERT FROM SINGLE TO DOUBLE DIMENSION
10     100 IJ=I*J+1
      NM=N*J+1
      DO 110 K=1,J
      NM=NM-NI
15     DO 110 L=1,I
      IJ=IJ-1
      NM=NM-1
      110 O(NM)=S(IJ)
      GO TO 140
      C CONVERT FROM DOUBLE TO SINGLE
20     120 IJ=0
      NM=0
      DO 130 K=1,J
      DO 125 L=1,I
      IJ=IJ+1
      NM=NM+1
25     125 S(IJ)=D(NM)
      130 NM=NM+NI
      140 RETURN
      10 END

```

SUBROUTINE MATPRO 74/74 OPT=1

FTN 4.2+75067 05/16/75 10.00.00.

```

      SUBROUTINE MATPRO(LOEM,A,B,PROD)
      C
      C THIS ROUTINE CALCULATES THE PRODUCT OF TWO MATRICES
      C
5     DIMENSION A(10,10),B(10,10),PROD(20,20)
      DOUBLE PRECISION A,B,PROD
      DO 20 I=1,LOEM
      DO 20 J=1,LOEM
      PROD(I,J)=0.
10     20 DO 20 L=1,LOEM
      PROD(I,J)=PROD(I,J)+A(I,L)*B(L,J)
      RETURN
      END

```

SUBROUTINE WRITE1 74/74 OPT=1

FTN 4.2+75067 05/16/75 10.01.12.

```

      SUBROUTINE WRITE1 (IQ,A)
      C
      C THIS ROUTINE WRITES A MATRIX
      C
5     DIMENSION A(10,10)
      DOUBLE PRECISION A
      DO 1 I=1,10
      1 WRITE(6,2) (A(I,J),J=1,10)
      2 FORMAT(6(4X,D16.9))
10     RETURN
      END

```

Figure 5. (Continued).

```

      SUBROUTINE VAND(REALP,COMPP,N,M,K,V)
      C
      C THIS ROUTINE CALCULATES THE VANDERMONDE FOR MATRICES WITH COMPLEX
      C EIGENVALUES
5      C
      DIMENSION REALP(1),COMPP(1),M(1),V(20,20)
      DOUBLE PRECISION V,REALP,COMPP
      DOUBLE PRECISION A,B,C,D,E,F,AA,BB
      NSUM=0
10      DO 10 L=1,K
          MM=M(L)
          IF (L.EQ.1) NC=1
          IF (L.GT.1) NC=NSUM+1
          NSUM=NSUM+MM
15      A=REALP(L)
          B=COMPP(L)
          C=1.00
          D=0.00
          DO 30 I=2,N
20      CALL COMCAL(A,B,C,D,E,F)
          C=E
          D=F
          V(I+N,NC)=D
          V(I,NC+N)=-C
25      V(I+N,NC+N)=V(I,NC)=C
30      CONTINUE
          V(1+N,NC+N)=V(1,NC)=1.00
          V(1+N,NC)=V(1,NC+N)=0.00
          IF (MM.EQ.1) GO TO 10
30      DO 40 I=1,N
          DO 40 J=2,MM
          JJ=NC+J-1
          IF (I-J) 97,98,99
          97 V(I+N,JJ+N)=V(I,JJ)=0.00
          V(I+N,JJ)=V(I,JJ+N)=0.00
          GO TO 40
          98 V(I+N,JJ+N)=V(I,JJ)=1.00
          V(I+N,JJ)=V(I,JJ+N)=0.00
          GO TO 40
          99 AA=V(I-1,JJ)
          BB=V(I+N-1,JJ)
          C=V(I-1,JJ-1)
          D=V(I+N-1,JJ-1)
          CALL COMCAL(AA,BB,A,B,E,F)
          E=E+C
          F=F+D
          V(I+N,JJ+N)=V(I,JJ)=E
          V(I+N,JJ)=F
          V(I,N+JJ)=-F
40      CONTINUE
50      CONTINUE
      RETURN
      END

```

Figure 5. (Continued).

SUBROUTINE COMCAL 74/74 OPT=1

FTN 4.2+75067

05/16/75 10.02.01.

```

      SUBROUTINE COMCAL(A,B,C,D,E,F)
      C
      C THIS ROUTINE CALCULATES THE PRODUCT OF A DOUBLE PRECISION COMPLEX
      C NUMBER
5      C
      C DOUBLE PRECISION A,B,C,D,E,F
      E=A*C-B*D
      F=A*D+B*C
      RETURN
10     ENO

```

SUBROUTINE VANDR 74/74 OPT=1

FTN 4.2+75067

05/16/75 10.02.53.

```

      SUBROUTINE VANDR(RELL,N,M,K,V)
      C
      C THIS ROUTINE CALCULATES THE VANDERMONDE FOR A MATRIX WITH ONLY REAL
      C EIGENVALUES
5      C
      C DIMENSION RELL(I),M(I),V(20,20)
      C DOUBLE PRECISION RELL,V
      C DOUBLE PRECISION Z,ZZ
      C NSUM=0
      C DO 10 L=1,K
      C MM=M(L)
      C IF(L.EQ.1) NC=1
      C IF(L.GT.1) NC=NSUM+1
      C NSUM=NSUM+MM
      C ZZ=1.00
      C Z=RELL(L)
      C DO 30 I=2,N
      C ZZ=ZZ*Z
      C V(I,NC)=ZZ
20      C 30 CONTINUE
      C V(I,NC)=1.00
      C IF(MM.EQ.1) GO TO 10
      C DO 40 I=1,N
      C DO 40 J=2,MM
      C JJ=NC+J-1
      C IF(I-J) 97,98,99
      C 97 V(I,JJ)=0.00
      C GO TO 40
      C 98 V(I,JJ)=1.00
      C GO TO 40
      C 99 V(I,JJ)=V(I-I,JJ-I)+V(I-I,JJ)*Z
      C 40 CONTINUE
      C 10 CONTINUE
      C RETURN
      C ENO

```

SUBROUTINE WRITER 74/74 OPT=1

FTN 4.2+75067

05/16/75 10.02.58.

```

      SUBROUTINE WRITER(XCOF,N)
      C
      C THIS ROUTINE WRITES A POLYNOMIAL
      C
5      C
      C DIMENSION XCOF(I)
      C DOUBLE PRECISION XCOF
      C WRITE(6,103) XCOF(1)
      C NN=N-1
      C DO 10 I=1,NN
      C 10 WRITE(6,100) XCOF(I+1),I
      C WRITE(6,101) XCOF(N+1),N
      C 100 FORMAT(/3X,1H(,ID15.8,1H),1X,4H*X**,12,1X,1H*)
      C 101 FORMAT(/3X,1H(,ID15.8,1H),1X,4H*X**,12)
      C 103 FORMAT(/3X,1H(,ID15.8,1H),1X,1H*)
      C RETURN
      C ENO

```

Figure 5. (Continued).

```

C
C
C
SUBROUTINE SORT(ROOTR,ROOTI,RELL,CMPXX,STORX,STORY,M,N,K,OK)
THIS ROUTINE SORTS THE ROOTS
5  DOUBLE PRECISION OK
   DIMENSION M(1)
   DIMENSION STORX(1),STORY(1)
   DIMENSION ROOTR(1),ROOTI(1),RELL(1),CMPXX(1)
10  DOUBLE PRECISION ROOTR,ROOTI,RELL,CMPXX
   DOUBLE PRECISION STORX,STORY
   DOUBLE PRECISION X,Y,W,Z,OO,DB
   K=0
   KK=0
15  DO 10 I=1,N
   K=K+1
   M(K)=0
   X=ROOTR(1)
   Y=ROOTI(1)
   RELL(K)=X
20  CMPXX(K)=Y
   NN=N-KK
   NK=0
   DO 20 J=1,NN
   W=ROOTR(J)
25  Z=ROOTI(J)
   OO=DABS(X-W)
   OB=DABS(Y-Z)
   IF (OO.LT.OK.AND.OB.LT.OK) GO TO 30
   NK=NK+1
30  STORX(NK)=W
   STORY(NK)=Z
   GO TO 20
40  M(K)=M(K)+1
20  CONTINUE
15  DO 40 L=1,NK
   ROOTR(L)=STORX(L)
   ROOTI(L)=STORY(L)
40  CONTINUE
   KK=KK+M(K)
   IF (KK.EQ.N) GO TO 50
10  CONTINUE
50  RETURN
END

```

Figure 5. (Continued).

```

      SUBROUTINE POLRT(XCOF,COF,M,ROOTR,ROOTI,IER)
      C
      C THIS ROUTINE CALCULATES THE ROOTS OF A POLYNOMIAL
      C
5     DOUBLE PRECISION FI
      DOUBLE PRECISION XCOF,COF,ROOTR,ROOTI
      DOUBLE PRECISION XO,YO,X,Y,XPR,YPR,UX,UY,V,YT,XT,U,XT2,YT2,SUMSQ
      DOUBLE PRECISION OX,OY,TEMP,ALPHA
      DIMENSION XCOF(1),COF(1),ROOTR(1),ROOTI(1)
10    IF(I1=0)
      N=M
      IER=0
      IF(XCOF(N+1)) 10,25,10
15    IF(N) 15,15,32
      IER=1
      RETURN
      IER=4
      GO TO 20
20    IER=2
      GO TO 20
32  IF(N=36) 35,35,30
      NX=N
      NXX=N+1
      N2=1
25    KJI=N+1
      DO 40 L=1,KJI
      MT=KJI-L+1
      COF(MT)=XCOF(L)
40    XO=.00500101
      YO=.01000101
30    IN=0
      X=XO
      XO=-10.0*YO
      YO=-10.0*X
15    X=XO
      Y=YO
      IN=IN+1
      GO TO 59
55  IF(I1=1)
      XPR=X
      YPR=Y
40    ICT=0
      UX=0.0
      UY=0.0
      V=0.0
      YT=0.0
      XT=1.0
      U=COF(N+1)
      IF(U) 65,130,65
50    DO 70 I=1,N
      L=N-I+1
      TEMP=COF(L)
      XT2=X*XT-Y*YT
      YT2=X*YT+Y*XT
      U=U+TEMP*XT2
      V=V+TEMP*YT2
      FI=1
55

```

Figure 5. (Continued).

SUBROUTINE POLRT 74/74 OPT=1

FTN 4.2*75067 05/16/75 10.08.24.

```

      UX=UX+FI*XT*TEMP
      UY=UY-FI*YT*TEMP
40      XT=XT2
      YT=YT2
      SUMSQ=UX*UX+UY*UY
      IF (SUMSQ) 75,110,75
      75      OX=(V*UY-U*UX)/SUMSQ
      X=X*OX
      65      UY=-(U*UY+V*UX)/SUMSQ
      Y=Y*OY
      78      IF (OABS(OY)+DABS(OX)-1.0-12) 100,80,80
      80      ICT=ICT+1
      70      IF (ICT-500) 60,85,85
      85      IF (IFIT) 100,90,100
      90      IF (IN-5) 50,95,95
      95      IER=3
      GO TO 20
      75      100      DO 105 L=1,NXX
      MT=XJ1-L+1
      TEMP=XCOF(MT)
      XCOF(MT)=COF(L)
      105      COF(L)=TEMP
      ITEMP=N
      NX=NX
      NX=ITEMP
      IF (IFIT) 120,55,120
      110      IF (IFIT) 115,50,115
      115      X=XPR
      Y=YPR
      120      IFIT=0
      122      IF (OABS(Y)-1.0-10=OABS(X)) 135,125,125
      125      ALPHA=X*X
      30      SUMSQ=X*X+Y*Y
      N=N-2
      GO TO 140
      130      X=0.0
      NX=NX-1
      35      NXX=NXX-1
      135      Y=0.0
      SUMSQ=0.0
      ALPHA=X
      N=N-1
      100      140      COF(2)=COF(2)+ALPHA*COF(1)
      145      DO 150 L=2,N
      150      COF(L+1)=COF(L+1)+ALPHA*COF(L)-SUMSQ*COF(L-1)
      155      ROOT1(N2)=Y
      ROOTR(N2)=X
      105      N2=N2+1
      IF (SUMSQ) 160,165,160
      160      Y=-Y
      SUMSQ=0.0
      GO TO 155
      110      165      IF (N) 20,20,45
      ENO

```

SUBROUTINE TRACE 74/74 OPT=1

FTN 4.2*75067 05/16/75 10.10.24.

```

      SUBROUTINE TRACE(A,N,TR)
      C
      C      THIS ROUTINE CALCULATES THE TRACE OF A MATRIX
      C
      5      DIMENSION A(10,10)
      DOUBLE PRECISION SUM,A,TR
      SUM=0.00
      DO 10 I=1,N
      10      SUM=SUM+A(I,I)
      TR=SUM
      RETURN
      ENO

```

SUBROUTINE TRACE2 74/74 OPT=1

FTN 4.2*75067 05/16/75 10.10.29.

```

      SUBROUTINE TRACE2(A,N,TR)
      C
      C      THIS ROUTINE CALCULATES THE TRACE OF A MATRIX
      C
      5      DIMENSION A(20,20)
      DOUBLE PRECISION SUM,A,TR
      SUM=0.00
      DO 10 I=1,N
      10      SUM=SUM+A(I,I)
      TR=SUM
      RETURN
      ENO

```

Figure 5. (Concluded).